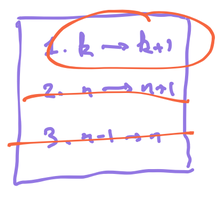
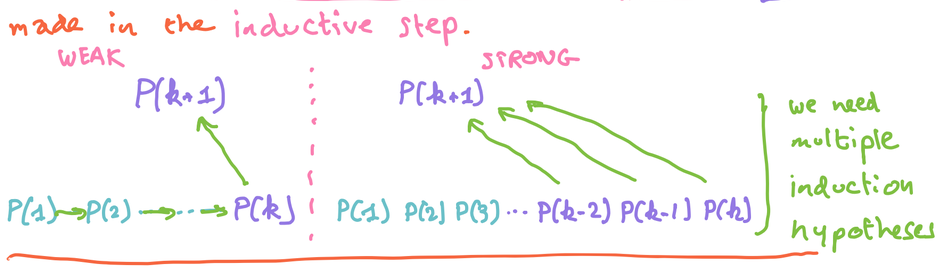
LECTURE 21: Strong induction continued



The only difference between weak and strong induction is the number of inductive hypotheses made in the inductive step.



Fundamental Theorem of Arithmetic Every integer n, $n>1$, can be written as a product of primes (unique up until order of the prime factors, i.e. $2\times7\times11 = 11\times2\times11 = \dots$)

Examples $24 = [4\times6] = 2\times2\times2\times3 = 2^3\cdot3^1$ ← $\cdot$ is multiplication

$100 = 2\times50 = 2\times2\times25 = 2\times2\times5\times5 = 2^2\cdot5^2$

$512 = 2^9$ ← product of primes

Proof Let $P(n)$: n can be written as a product of primes

BASE case: $n=2$ we can verify that $2 = 2^1$ and since 2 is a prime number 2^1 is a valid product of primes. ✓

INDUCTIVE case: ⚠ If we were doing weak induction we would assume $P(k)$ is true, and show $P(k+1)$

Instead: We assume that $P(i)$ is true for $2\le i\le k$, in other words: every number smaller or equal to k can be decomposed.

BASE CASE $2\le i\le k$

STRONG INDUCTION HYPOTHESIS

Let's consider the integer $k+1$.

[Because we are interested in decomposing $k+1$ as a product of primes, then it is normal to ask first: is $k+1$ prime or not?]

- **CASE 1:** $k+1$ is PRIME, then it is its own decomposition, therefore $P(k+1)$ is TRUE ✓
- **CASE 2:** $k+1$ is not prime, it is COMPOSITE

  By DEFINITION of "COMPOSITE" this means there exists

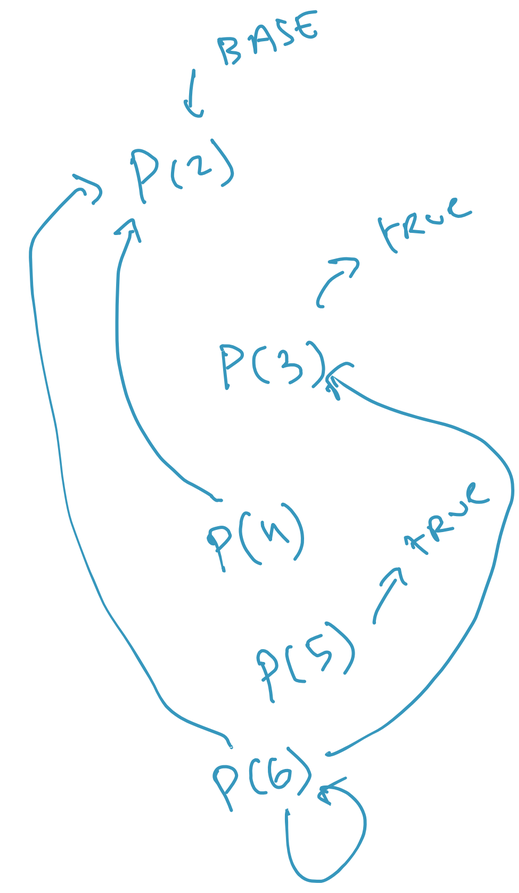
  $\exists a,b\in\mathbb{Z},\ a>1,\ b>1,\ k+1 = a\cdot b$

  Because $a<k+1$ and $b<k+1$, by induction hypothesis, there is a prime factor decomposition for both a and b, and therefore a prime decomposition for $a\cdot b = k+1$.

Therefore we have shown that $P(k+1)$ is true.

CONCLUSION: We can conclude by strong induction that $P(n)$ is true for all $n\ge2$, $n>1$

What we know to be true to think of strong induction →



How to translate COMPOSITE to definition?

- $x\in A\setminus B$ ↔ $x\in A$ AND $x\notin B$
- n is even ↔ $\exists k\in\mathbb{Z}$, $n=2k$
- $\exists a,b\in\mathbb{Z}$, $a>1$, $b>1$, $n = a\cdot b$

→ without this condition then every integer (including prime numbers) would be considered "composite"

$7 \rightsquigarrow 1\times7\times1$

Binary Decomposition Theorem

Every number n can be written in binary form

$$n = b_k\cdot2^k + b_{k-1}2^{k-1} + \dots + b_2 2^2 + b_1 2^1 + b_0 2^0$$

b_k is the k-th bit, $b_k = 0$ or $b_k = 1$

Examples: $4 = 1\cdot2^2$ $\quad 7 = 1\cdot2^2 + 1\cdot2^1 + 1\cdot2^0$

PROOF.

WEAK INDUCTION

$P(n)$: n can be written (uniquely) in binary

BASE CASE: $n=1$, for $n=1$, we have the decomposition $1 = 1\cdot2^0$ ✓

INDUCTIVE case: We assume that $P(k)$ is true and we want to show $P(k+1)$ is true.

CASE 1 is similar as strong induction, because in the strong induction we only rely on the previous number.

CASE 2 is much more complicated because you have to transition from k to $k+1$.

Hint: consider that you are adding two numbers with a CARRY

```
  1 1 1  ← CARRIES
  1 0 0 x x x
  0 0 0 0 0 x
-------------
  1 0 1 0 0 0
```

STRONG INDUCTION

$P(n)$: n can be written (uniquely) in binary

BASE CASE: $n=1$, for $n=1$, we have the decomposition $1 = 1\cdot2^0$ ✓

INDUCTIVE case: We assume that for all i, $1\le i\le k$, $P(i)$ is true.

We want to show that $P(k+1)$ is true.

CASE 1 $k+1$ is odd. Therefore it can be expressed as

$k+1 = k + 1\cdot2^0$ ($1\cdot2^0$ is just a fancy way of rewriting 1)

We can then, by inductive hypothesis, decompose k, and then set the lowest bit to 1.

CASE 2 $k+1$ is even. Therefore it can be expressed as $k+1 = 2p$ (for $p\in\mathbb{Z}$) and because $p<k+1$, we can apply the inductive hypothesis on p to get a binary decomposition. When multiplying we are shifting the bits by one base

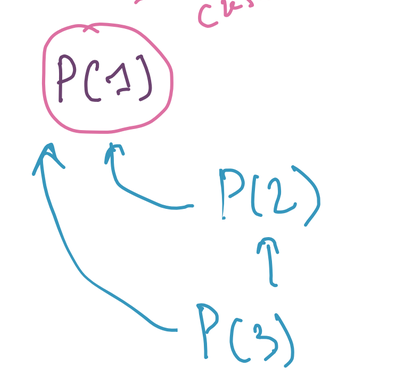
```
  5      1 0 0
2×5      1 0 1 0
```

CONCLUSION: We have shown by strong induction $P(n)$ is true

some binary numbers	
0	000 ← even
1	001 ← odd
2	010 ← even
3	011 ← odd
4	100 ←
5	101 ← odd

$\forall n,\ P(n)$

$\forall n\ge5,\ P(n)$



don't care about p

$k+1 = 2p+1$

Q details

$k+1$ is odd

therefore

k is even

$k = \dots\dots0$ ← last bit

\# because k is smaller, by inductive hypothesis we can decompose it

* because k is even, its last bit is 0

* FINALLY to combine the decompositions of k and 1 to get $k+1$, we simply take the decomposition of k and set the lowest bit of this to 1 to get a decomposition of $k+1$