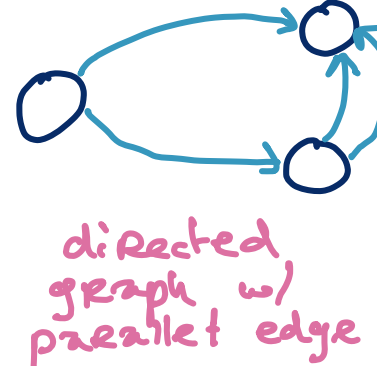
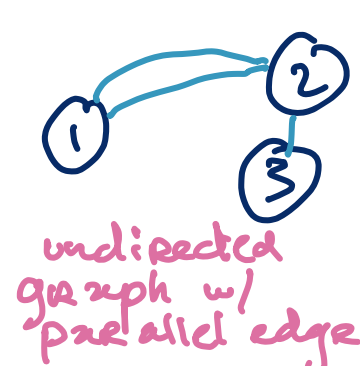


## Lecture 24

### Quick answers

- parallel edges can be directed or undirected



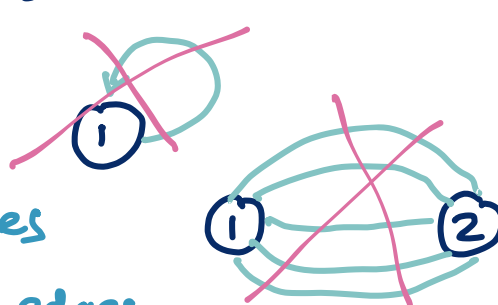
- not parallel edge if different direction



not a parallel edge  
 $(1, 2) \neq (2, 1)$   
 $(A, B) \neq (B, A)$

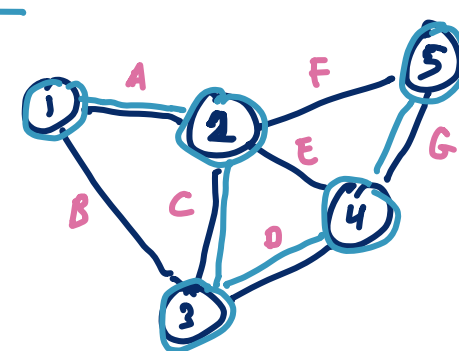
In this course, unless explicitly mentioned, we are interested in graphs that are:

- undirected
- no self-loops
- no parallel edges
- finite number of edges (and vertices)



### PROPERTIES OF GRAPHS

Def.: A path is an alternating sequence of vertices and edges, starting and ending in a vertex, and no vertex appears more than once.

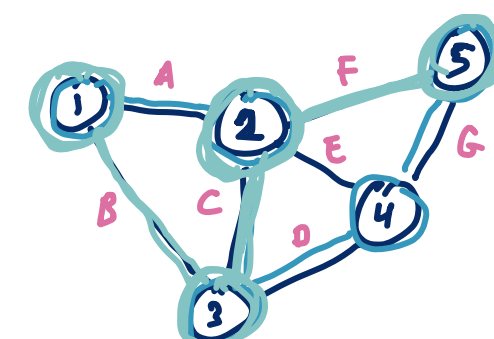


this is an alternating sequence

1 A 2 C 3 D 4 G 5  
 (example of a path from 1 to 5)

Def.: A walk is

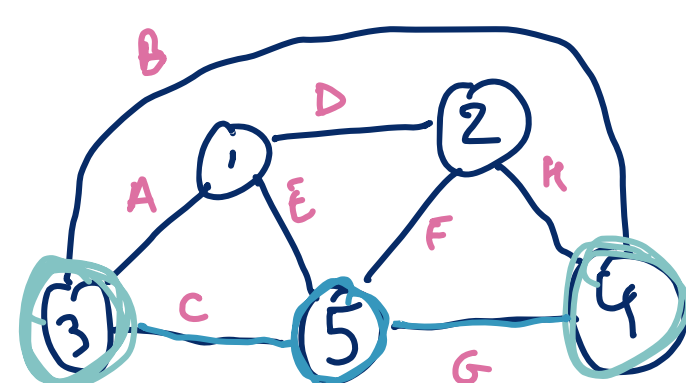
- an alternating seq. of vertices and edges,
- starting and ending in a vertex.



therefore a path cannot contain a loop

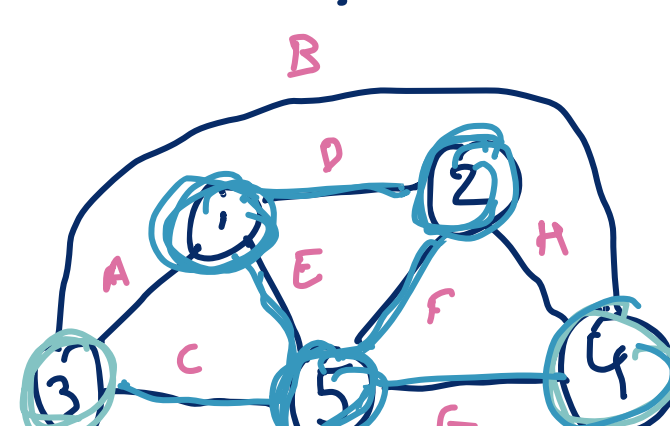
1 B 3 C 2 F 5  
 (another path 1 to 5)

Question: What is the difference between a path and a walk?



3 C 5 G 4

PATH = no repetition



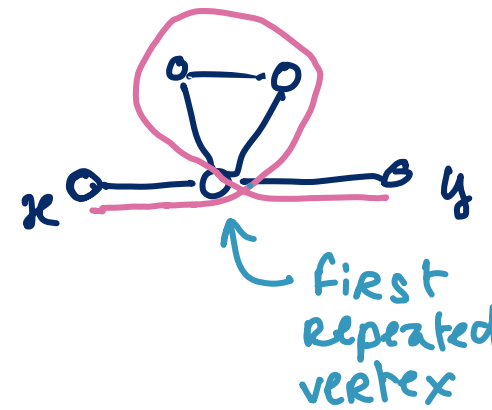
3 C 5 E 1 D 2 F 5 E 1 D 2 F 5 E 1 D 2 F 5 E 1 D 2 F 5 G 4

WALK = allows for repetition

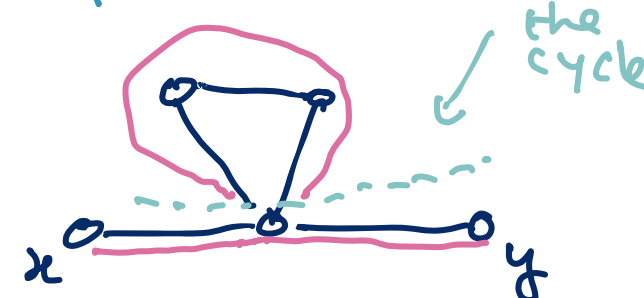
Corollary: Every path is a walk.

Thm: If there is a WALK between two vertices  $x$  and  $y$ , there must be a PATH between  $x$  and  $y$ .

If you have a walk:



you can simply "cut" the cycle and you get a path:



Def: A cycle is a walk that starts and ends in the same vertex.

Begin by considering any walk between vertices  $x$  and  $y$ . This walk can be categorized in two cases:

CASE 1 The walk does not contain repeated vertices.

- in this case, the walk already qualifies as a path, by definition.

CASE 2 The walk includes at least one vertex that is repeated.

- Let's call  $z$  the first repeated vertex.

$w = (\dots z, e_1, \dots, e_k, z, \dots)$

- Create a shorter walk by eliminating the segment of the walk between the first and last occurrences of  $z$ .

$w' = (\dots z \dots)$

Specifically, we only retain the first occurrence of  $z$ . The resulting is strictly shorter, because we have removed at least one vertex [Argument to avoid infinity.]

- Continue this process iteratively until there are no more repeated vertices.
- Since each iteration removes at least one vertex, and there are a finite number of vertices, this process will eventually result - in a walk with no repeated vertices - that is, a path!

### CONCLUSION:

- The final walk goes from  $x$  to  $y$  without repeating vertices, thus meeting the def. of a path
- This process proves that if there is a walk between two vertices, there must also be a PATH.

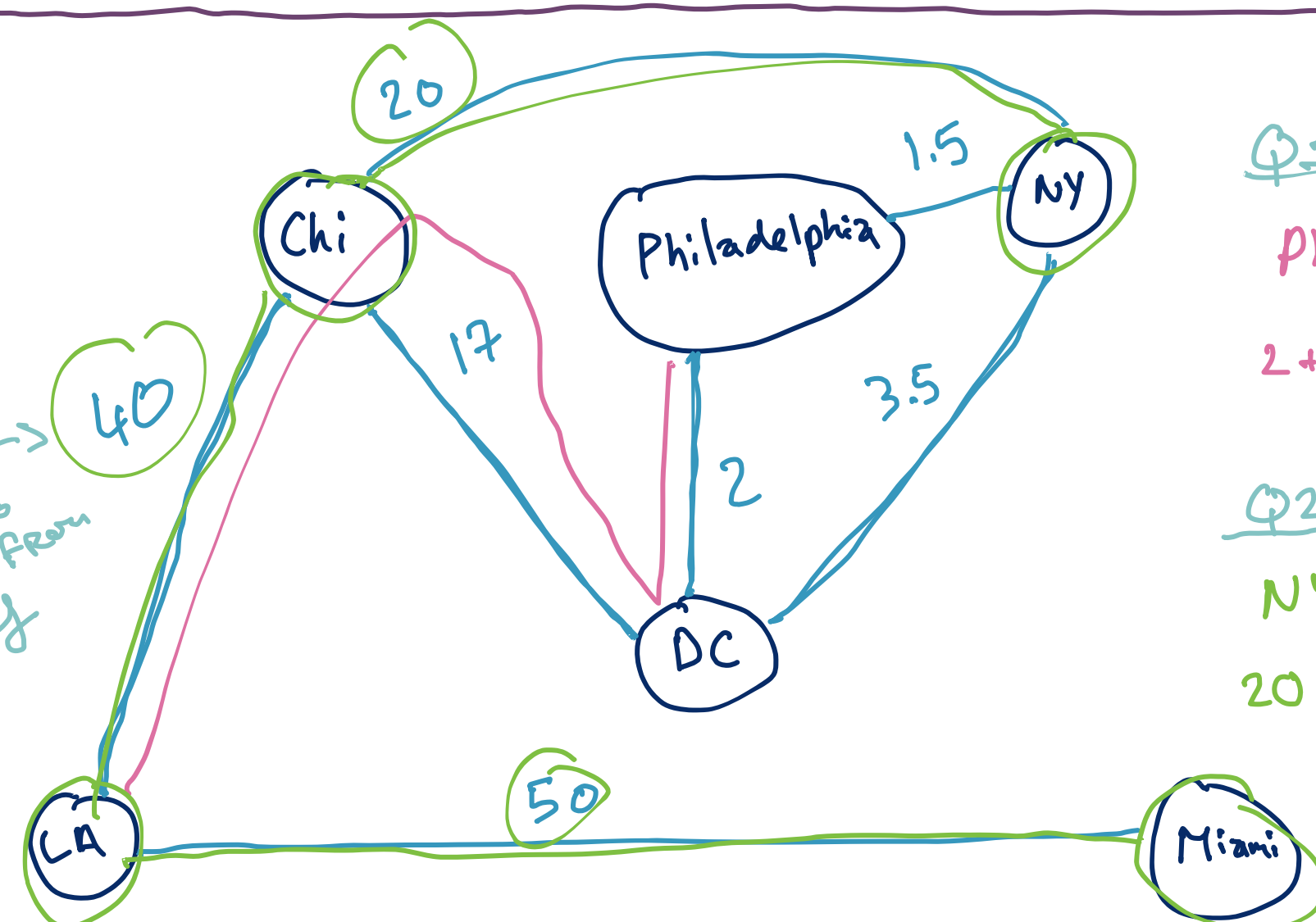
formal notation to represent PATH / WALK / CYCLE

$e_i$  this is an edge

$w = (v_1, e_1, v_2, e_2, v_3, e_3 \dots)$

this is vertex  $v_i$

### MODELING THE RAIL NETWORK WITH GRAPHS



Q1: Phil  $\rightarrow$  LA?

Phil  $\rightarrow$  DC  $\rightarrow$  Chi  $\rightarrow$  LA

$2 + 17 + 40 = 59$

Q2: NY  $\rightarrow$  Miami

NY  $\rightarrow$  Chi  $\rightarrow$  LA  $\rightarrow$  Miami

$20 + 40 + 50 = 110$